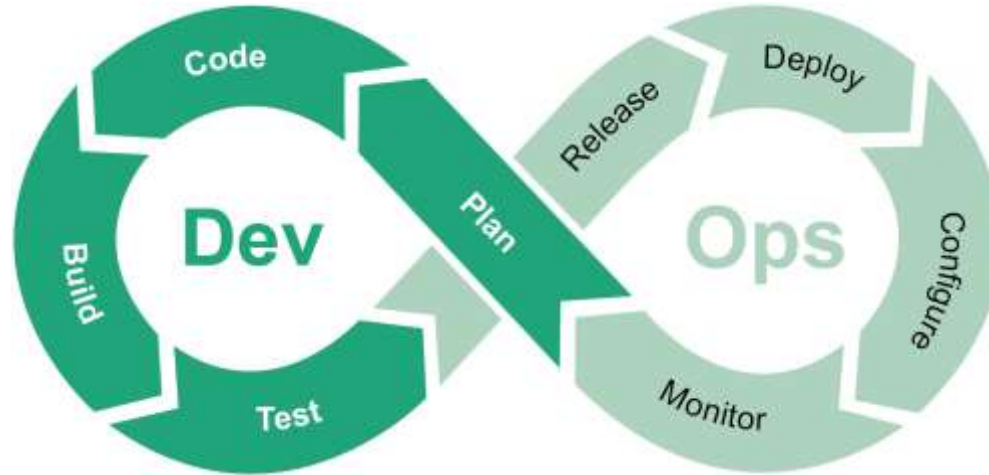


---

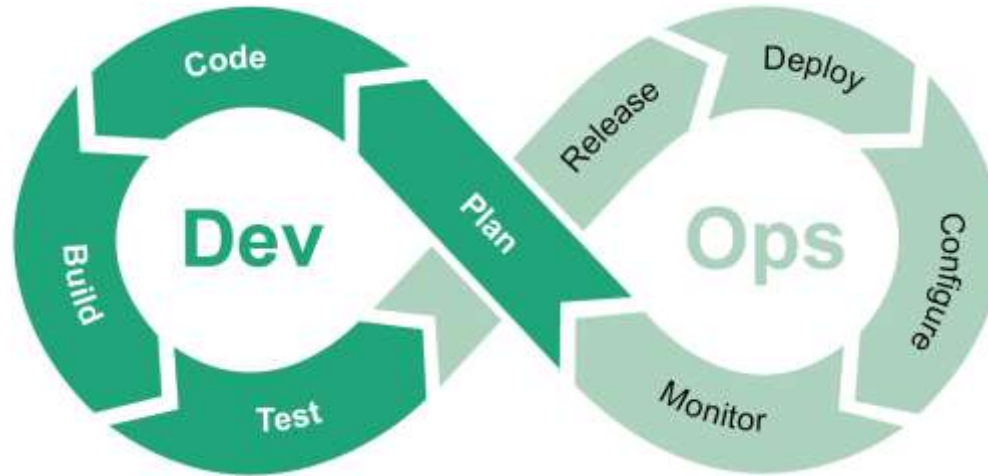
# ENABLING DEVSECOPS WITH THE NEW GENERATION OF STATIC ANALYSIS

Prof. Dr. Eric Bodden

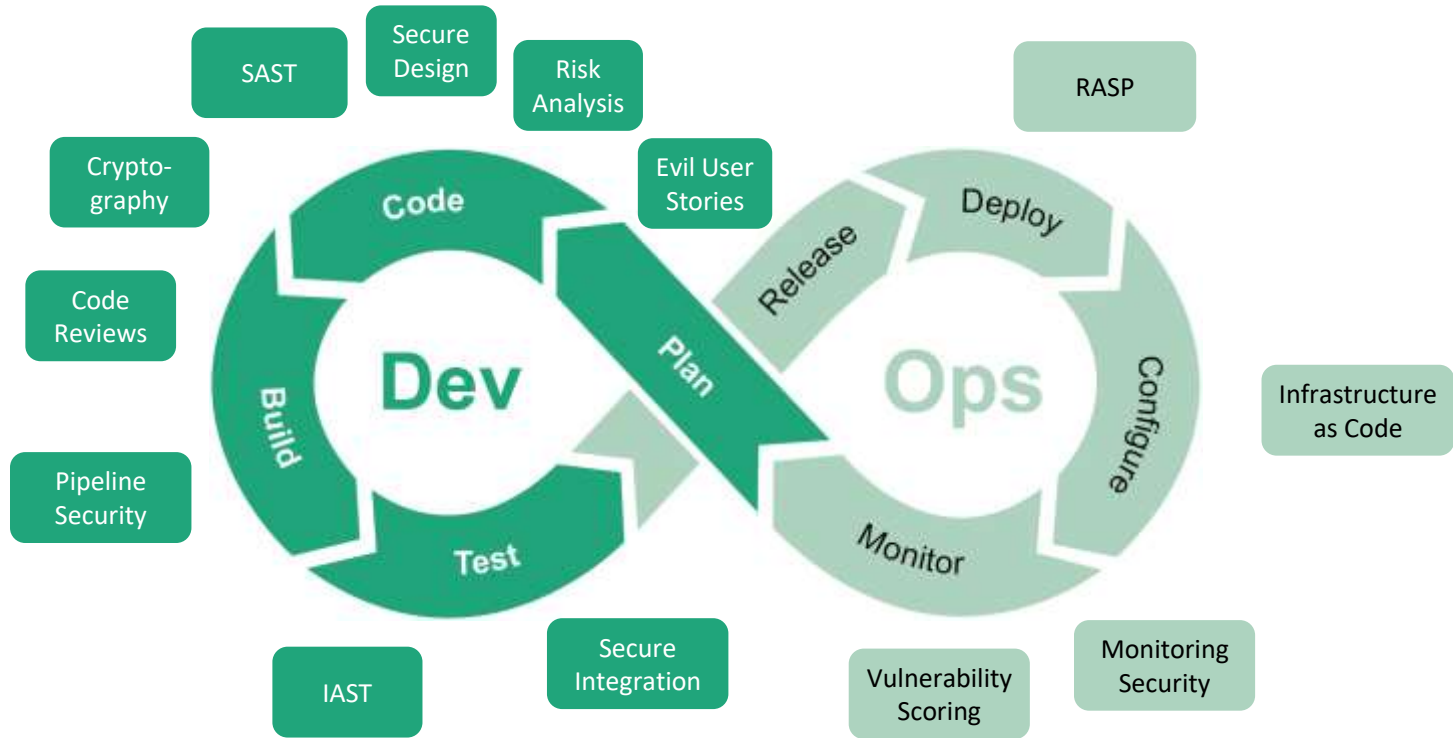
---



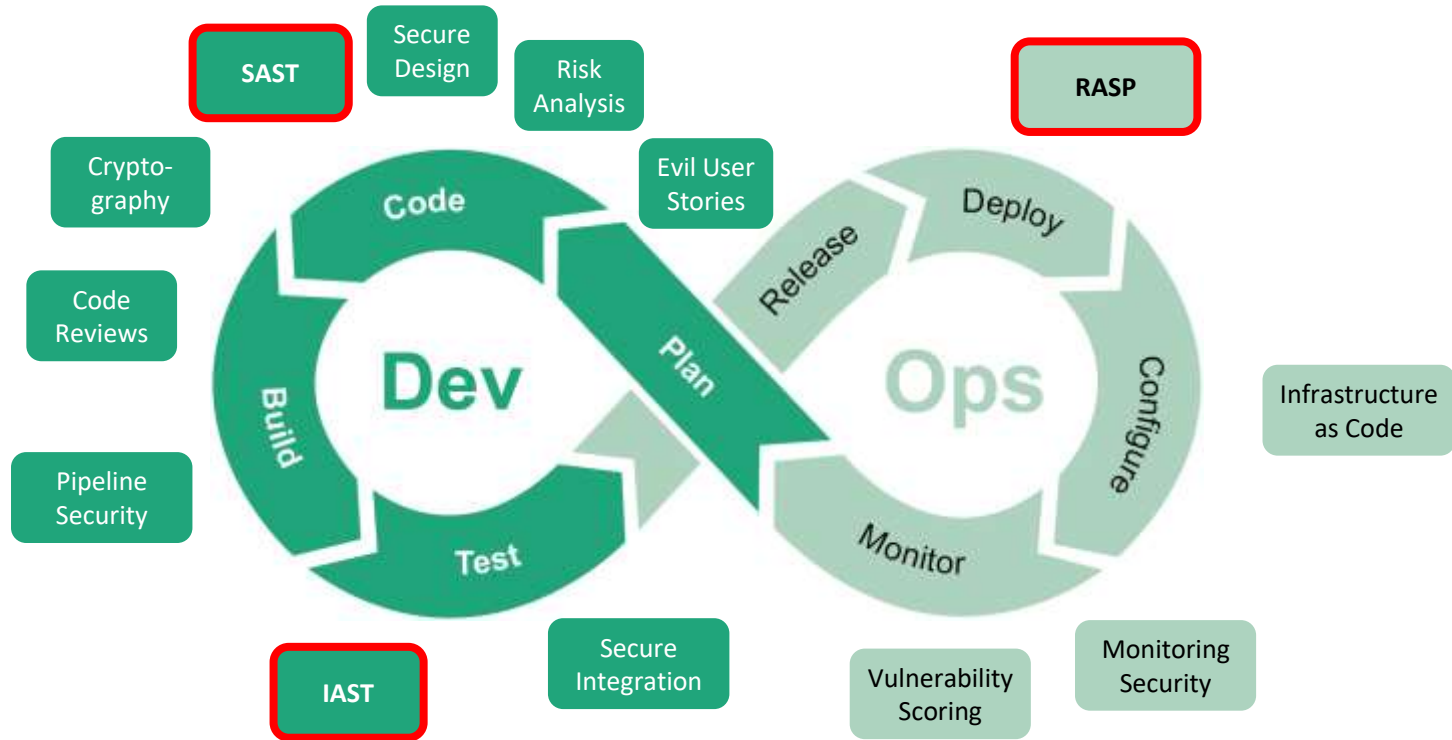
# DevOps



# Security touchpoints in DevSecOps



# Focus here: tool automation for security testing



# SAST vs. IAST vs. RASP

## Static Application Security Testing

- Analyzes program code without executing it
- Can cover entire program, not just a single run
- Works on incomplete code
- White-box technique:
  - Can give immediate feedback to developer
  - Can instantly suggest sensible fixes
- But: traditionally many false warnings

## Interactive Application Security Testing

- Analyzes program code during execution, alerting the developer about detected vulnerabilities
- Requires executable code
- Often too expensive for deployment / ops
- Good: Can be very precise, little to no false warnings

## Runtime Application Self-Protection

- Analyzes program code during execution, automatically mitigating vulnerabilities
- Has to be efficient, thus often restricted to efficiently recognizable properties
- Sensible to use as an additional layer of defense

# At Fraunhofer IEM we focus on SAST, i.e., static analysis

## Static Application Security Testing

- Analyzes program code without executing it
- Can cover entire program, not just a single run
- Works on incomplete code
- White-box technique:
  - Can give immediate feedback to developer
  - Can instantly suggest sensible fixes
- But: traditionally many false warnings

# At Fraunhofer IEM we focus on SAST, i.e., static analysis

## Static Application Security Testing

- Analyzes program code without executing it
- Can cover entire program, not just a single run
- Works on incomplete code
- White-box technique:
  - Can give immediate feedback to developer
  - Can instantly suggest sensible fixes
- **But: traditionally many false warnings**

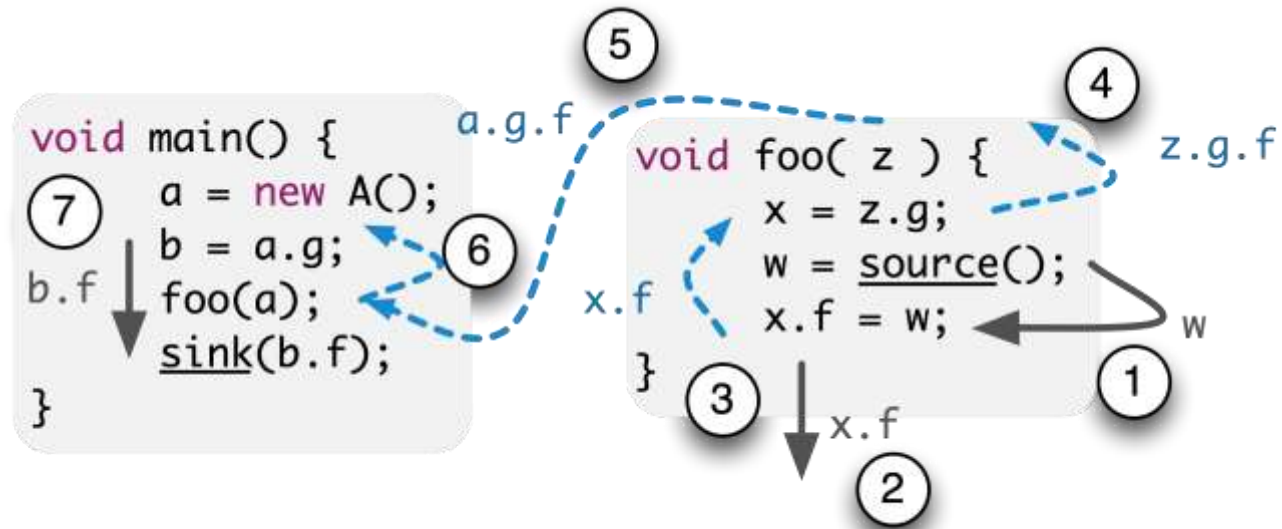
Why is precise static analysis so challenging?

# 1st Major challenge: Heap assignments and “Aliasing”

```
String secret = secret();
```

```
print(output);
```





Will it leak?

## 2nd Major challenge: Field-sensitivity

```
a.f = mySecret();  
print(a.g);
```

can represent a.f as:

a.\*

field-insensitive

## 2nd Major challenge: Field-sensitivity

```
a.f = mySecret();  
print(a.g);
```

False positives because  
analysis assumes also a.g as tainted

a.\*

field-insensitive

## 2nd Major challenge: Field-sensitivity

```
a.f = mySecret();  
print(a.g);
```

can represent a.f as:

a.\*

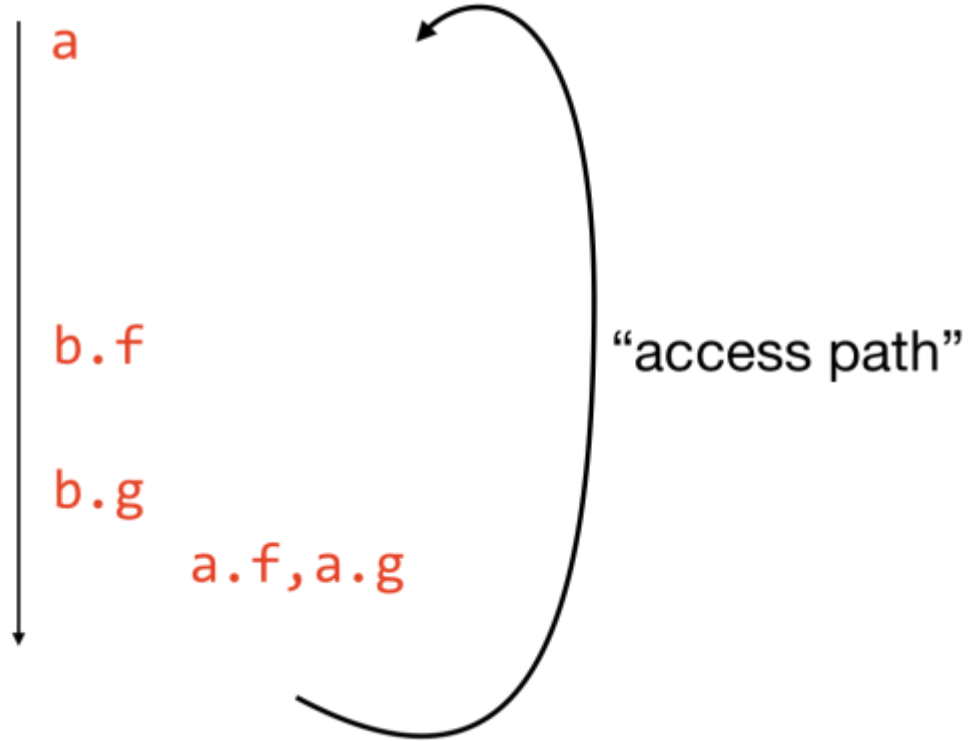
a.f

field-insensitive

field-sensitive

## 2nd Major challenge: Field-sensitivity

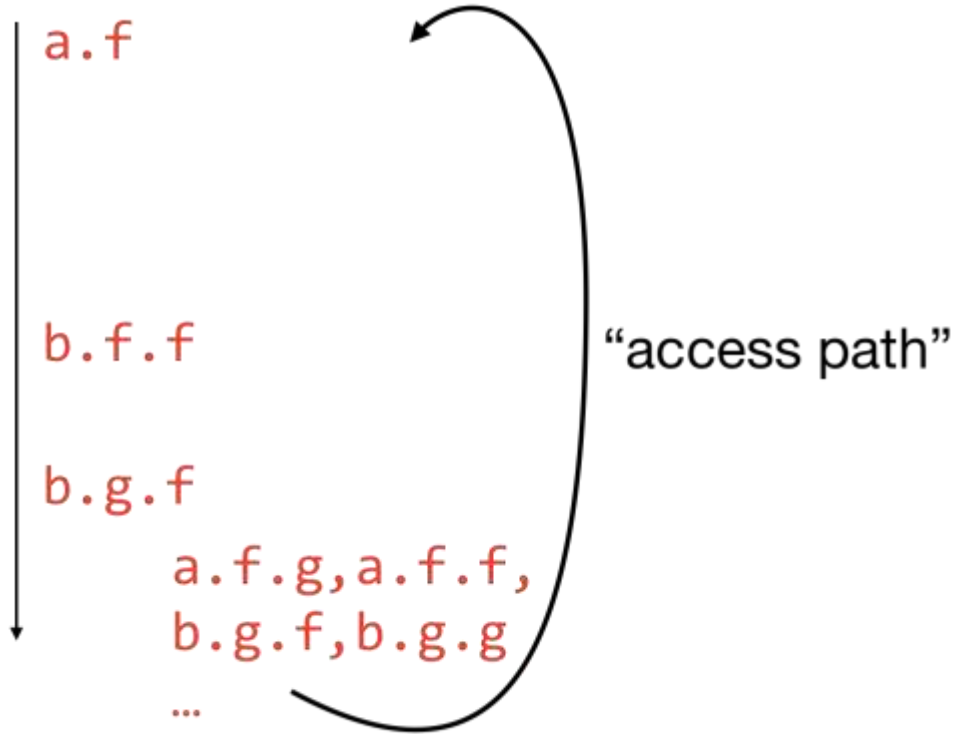
```
a = mySecret();  
while(..) {  
  A b = new A();  
  if(..)  
    b.f = a;  
  else  
    b.g = a;  
  a = b;  
  b = null;  
}
```



## 2nd Major challenge: Field-sensitivity

```
a = mySecret();  
while(..) {  
  A b = new A();  
  if(..)  
    b.f = a;  
  else  
    b.g = a;  
  a = b;  
  b = null;  
}
```

require finite representation!



Using novel algorithms we are now able to analyze program code with almost no loss of precision

In large-scale empirical evaluations: false-positive rates below 5%!

## Context-, Flow-, and Field-Sensitive Data-Flow Analysis using Synchronized Pushdown Systems

JOHANNES SPÄTH, Fraunhofer IEM, Germany  
KARIM ALL, University of Alberta, Canada  
ERIC BODDEN, Heinz Nixdorf Institut, Universität Paderborn and Fraunhofer IEM, Germany

Precise static analyses are context-, field- and flow-sensitive. Context- and field-sensitivity are both expressible in context free language (CFL) reachability problems. Solving both CFL problems along the same data-flow path is undecidable, which is why most flow-sensitive data-flow analyses over-approximate field-sensitivity through  $k$ -limited access paths, or through access graphs. Unfortunately, as our experience and this paper show, both representations do not scale very well when used in analyzer programs with recursive data structures.

Any single CFL-reachability problem is efficiently solvable, by means of a pushdown system. This work thus introduces the concept of *synchronized pushdown systems* (SPDS). SPDS encode both procedure calls/returns and field stores/loads as separate but “synchronized” CFL reachability problems. An SPDS solves both individual problems precisely, and approximation occurs only in corner cases that are apparently rare in practice: at statements where both problems are satisfied but not along the same data-flow path.

SPDS are also efficient: formal complexity analysis shows that SPDS shift the complexity from  $|P|^{14}$  under  $k$ -limiting to  $|P|^2$ , where  $P$  is the set of fields and  $S$  the set of statements involved in a data flow. Our evaluation using DaCapo shows this shift to pay off in practice: SPDS are almost as efficient as  $k$ -limiting with  $k = 1$  although their precision equals  $k = \infty$ . For a separate analysis SPDS accelerate the analysis up to 85% for data flows of objects that involve many field accesses but span rather few methods.

We conclude that SPDS can provide high precision and further improve scalability, in particular when used in analyses that expose rather local data flows.

CCS Concepts: Theory of computation → Program analysis; Object oriented concepts; Grammars and context free languages.

Additional Key Words and Phrases: static analysis, data-flow, aliasing, access paths, pushdown system

### ACM Reference Format:

Johannes Späth, Karim All, and Eric Bodden. 2019. Context-, Flow-, and Field-Sensitive Data-Flow Analysis using Synchronized Pushdown Systems. *Proc. ACM Program. Lang.* 3, POPL, Article 48 (January 2019), 29 pages. <https://doi.org/10.1145/3290361>

### 1 INTRODUCTION

Static data-flow analysis helps detect bugs and security vulnerabilities early in the software development process, including semantic properties such as null-pointers [Narda and Sieha 2009], data races [Gibson et al. 2009; Yan et al. 2011], and misuse of application programming interfaces

Authors' addresses: Johannes Späth, Fraunhofer IEM, Germany, johannes spaeth@ipm.fraunhofer.de; Karim All, University of Alberta, Canada, karim.all@ualberta.ca; Eric Bodden, Heinz Nixdorf Institut, Universität Paderborn and Fraunhofer IEM, Germany, eric.bodden@uni-paderborn.de

Permission to make digital or hard copies of all or part of this work for personal or classroom use is granted without fee provided that copies are not made or distributed for profit or commercial advantage and that copies bear this notice and the full citation on the first page. Copyrights for components of this work owned by others than ACM must be honored. Abstracting with credit is permitted. To copy otherwise, to republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee. Request permissions from [permissions.acm.org](http://permissions.acm.org).

© 2019 Association for Computing Machinery.  
1076-1421/2019/1-ART248  
<https://doi.org/10.1145/3290361>

*Proc. ACM Program. Lang.*, Vol. 3, No. POPL, Article 48. Publication date: January 2019.

48

[POPL 2019] ACM Distinguished Paper



available at <http://eclipse.org/cognicrypt/>

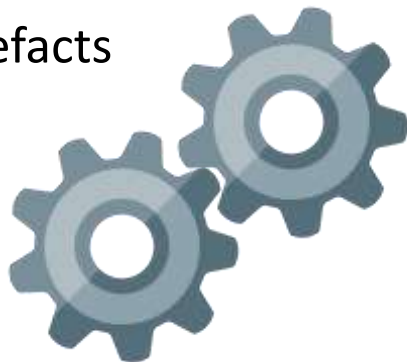


# Analysis of MavenCentral



2.7+

Million  
Software Artefacts



73%

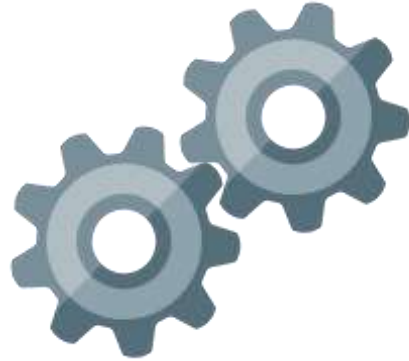
of all artefacts  
use crypto  
in an  
insecure  
manner!

[A Large-Scale Study of Non-Trivial Misuses of the Java Cryptography Architecture. Stefan Krüger Johannes Spaeth, Karim Ali, Eric Bodden, Mira Mezini.]

# Analysis of Google Play: top banking/health/password apps

250

Apps



71%

of all apps  
use crypto  
in an  
insecure  
manner!

[A Large-Scale Study of Non-Trivial Misuses of the Java Cryptography Architecture. Stefan Krüger Johannes Spaeth, Karim Ali, Eric Bodden, Mira Mezini.]



# Norton Identity Safe Password

NortonMobile Tools

★★★★★ 20,228

USK: All ages

This app is compatible with all of your devices.

Add to Wishlist

Install



One of many examples:

Norton Identity Safe

> 500.000 Downloads in Google Play

# Norton Identity Safe Privilege Escalation

SYMSA1460 | August 29th, 2018 | <https://www.symantec.com/docs/SYMSA1460>

Security Advisory ID	SYMSA1460
Initial Publication Date:	29 Aug 2018
Advisory Status:	Closed
Advisory Severity:	Medium
CVSS Base Score:	5.6

CVE-2018-12240

## Summary

Symantec has released an update to address an issue that was discovered in the Norton Identity Safe for Android product.

★	Was this article helpful?	
	No	Yes
	Print Article	
	Subscribe to this Article	
	Manage your Subscriptions	



# VR-Banking

Fiducia & GAD IT AG Finanzen

★★★★★ 19.702

USK ab 0 Jahren

Diese App ist mit allen deinen Geräten kompatibel.

Zur Wunschliste hinzufügen

## Zertifiziertes Produkt

Geprüfte Online-Applikation:

Zertifikatsinhaber: Fiducia & GAD IT AG

Prüfzeichennummer: 0000043889



Datenschutz/  
Datensicherheit

www.tuv.com  
ID 0000043889

Die Prüfung umfasst:

- Datenschutz/ Datensicherheit



## Informationen

### Beschreibung:

Für die Online-Banking-Applikationen eBanking Private Edition und eBanking Business Edition sowie die VR-BankingApp und alle auf deren Basis individualisierten Banking-Apps und Applikationen (siehe Service Information) hat die Fiducia & GAD IT AG einen wirksamen Prozess zur Erreichung folgender Ziele etabliert:

- Vertraulichkeit und Integrität der verarbeiteten Informationen
- Wirksame Umsetzung der Aussagen der Datenschutzerklärung
- Wirksamer Schutz der personenbezogenen Daten gemäß anwendbarer, aktueller Datenschutzgesetzgebung
- Wirksame Absicherung der von außen zugänglichen technischen Systeme gegen unbefugte Nutzung

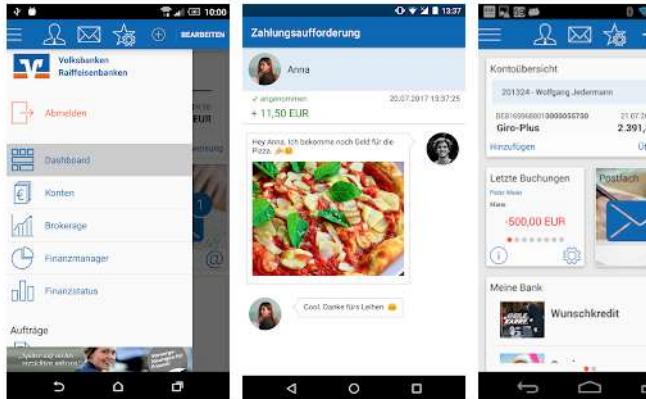
Zusätzlich wurde geprüft, ob die zur Autorisierung der Banking-Transaktionen genutzten Verfahren mobileTAN und Smart-TAN nach Best-Practices implementiert wurden.

Der Nachweis wurde durch ein Datenschutzaudit sowie externe und interne Sicherheitsanalysen erbracht. Der Prüfbericht Nr. 63008709-01 in der aktuellen Version ist Bestandteil dieses Zertifikats.

Die Wirksamkeit des geprüften Prozesses wird durch die TÜV Rheinland i-sec GmbH regelmäßig überwacht.

Das Zertifikat basiert auf einem von der TÜV Rheinland i-sec GmbH entwickelten Anforderungskatalog und stellt kein akkreditiertes Zertifizierungsverfahren, Siegel oder Prüfzeichen im Sinne der Art. 42, 43 der Verordnung (EU) 2016/679 (Datenschutz-Grundverordnung) dar.

Dieses Zertifikat ist gültig bis 13.12.2020.



Ihre Finanzen im Blick – immer und überall mit der VR-Banking App der Volksbank Raiffeisenbanken.

Sie möchten unterwegs mal eben den Kontostand abrufen, kurz die letzten Umsätze prüfen, eine dringende Überweisung erledigen, Börseninformationen einholen und kein Problem mit der VR-Banking App. Besonders praktisch: legen Sie Ihre beliebte

WEITERE INFORMATIONEN



# Integrates with your favourite CI-environment

quests Issues Marketplace Features

johspaeth /

Code

Create Conversation

Open

Conversation

johspaeth

Hey Ev

JCodeShield

```
CryptoUtils.java
```

```
private static SecretKey deriveSecretKey(char[] passPhrase) {  
    try {  
        PBEKeySpec keySpec = new PBEKeySpec(passPhrase, SALT, KDF_ITERATION_COUNT, ENCRYPTION_KEY_L
```

JCodeShield on 4 May • edited by johspaeth

JCodeShield found two issue in this line.

Usage of a hardcoded password via parameter `passPhrase` (value is `"elasticsearch-license".toCharArray()`) as defined in [line 39](#). [CWE-259](#)

The parameter `SALT` is not randomly generated (hardcoded in [line 32](#)). An attacker having access to the source code (or who can decompile the bytecode) can extract the `SALT`, derive the key from the password and break the encryption. [CWE-759](#)

JCodeShield on 4 May • edited by johspaeth

JCodeShield found four issues in this line.

The parameter `hashedPassphrase` is hardcoded to `"elasticsearch-license"` ([line 166](#)). [CWE-259](#).

Milestone: No milestone

Notifications: Customize

Unsubscribe

HEINZ NIXDORF INSTITUT  
UNIVERSITÄT PADERBORN

Fraunhofer IEM

# Visit us at Hall 10, Booth 311

```
public class Msg {  
    public byte[] sign(String data) throws InvalidKeyException, Exception {  
        return null;  
    }  
}
```

